

## Considerations for LS-DYNA Efficiency in SGI IRIX and Linux Environments with a NUMA System Architecture

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### Keywords:

Mechanical computer-aided engineering, high performance computing, finite element analysis, impact simulation

**ABSTRACT**

Manufacturing industry and research organizations continue to increase their investments in structural analysis and impact simulations such that the growing number of LS-DYNA users continue to demand more from computer system environments. These demands typically include rapid single job turnaround and multi-job throughput capability in a multi-user high-performance computing (HPC) environment.

What is more, many LS-DYNA simulation environments coexist with other mechanical computer aided engineering (MCAE) software for structural analysis and computational fluid dynamics that all compete for system resources such as CPU cycles, system bandwidth, memory, disk storage and I/O.

This paper examines the computational efficiency of structural analyses and simulations for relevant applications in LS-DYNA. Parameters such as model size, element types, and simulation conditions can produce a wide range of computational behavior such that consideration should be given to how system resources are allocated and configured.

The computational characteristics of the SGI® Origin® 3000 servers, based on IRIX® and MIPS®, and the SGI® Altix™ 3000 servers, based on Linux® and Itanium® 2 from Intel®, are examined for both turnaround and throughput requirements that include industrial-size examples. In addition, simple guidelines on proper system configuration and innovative use of available SGI system resource management tools are provided that are designed to maximize LS-DYNA productivity.

**Introduction**

Mechanical design and manufacturing organizations increasingly rely on high-performance computing (HPC) technology and mechanical computer-aided engineering (MCAE) applications to drive innovation in product development. Industries such as automotive, aerospace, turbomachinery, and others are challenged with an increasing need to reduce development cycles, while satisfying global regulations on safety, environmental impact and fuel efficiency. They must also appeal to demands for high quality, well-designed products in a competitive business environment.

Continuing technology advances in finite element analysis (FEA) provide industry with a design aid that is a relevant step towards achieving their product development goals. Historically, FEA simulation provided limited value as an influence on industrial design owing to excessive modeling and solution times that could not meet conventional development schedules. During the 1980's vector architectures offered greatly improved FEA simulation turn-around, but at a very high cost. Scalable RISC computing introduced in the 1990's narrowed this gap of cost-performance.

Recent advancements in parallel computing have demonstrated that vector-level performance can be easily exceeded with proper implementation of parallel FEA algorithms for distributed shared-memory servers like the SGI Origin 3000 and clusters like the SGI Altix 3000. Perhaps even more appealing is that the increased performance is offered at a fraction of the cost. These trends have influenced recent increased investments by users of FEA technology throughout a range of industries.

The HPC paradigm, however, has been undergoing a change in recent years thanks to a proliferation of Linux-based scalable systems. Such systems equipped with Intel

processors and some, such as the SGI Altix 3000 that offer an inexpensive yet innovative memory design, can deliver a level of performance that rivals vector and most RISC systems in nearly all aspects. For example, all conventional clusters are capable of processing thousands of MIPS and hundreds of MFLOPS, a level well beyond the single digit performance in the mid-1980's RISC when it was first introduced.

With release of LS-DYNA 940 during 1999, the first substantial parallel capability for explicit FEA simulation provided the industry with a migration path from expensive vector systems to more economical scalable systems. A distributed memory parallel (DMP) programming model is used for LS-DYNA in order to leverage the advantage of these contemporary scalable architectures. This paper examines efficient implementation of LS-DYNA for both uni-processor and parallel considerations, and for scalable RISC servers and 64-bit Linux clusters.

### **HPC Characteristics of LS-DYNA**

Finite element analysis software LS-DYNA™ from Livermore Software Technology Corporation ([www.lstc.com](http://www.lstc.com)) is a multi-purpose structural and fluid analysis software for high-transient, short duration structural dynamics, and other multi-physics applications. Considered one of the most advanced nonlinear finite element programs available today, LS-DYNA has proved an invaluable simulation tool for industry and research organizations who develop products for automotive, aerospace, power-generation, consumer products, and defense applications, among others.

LS-DYNA simulations for the automotive industry include vehicle crash and rollover, airbag deployment and occupant response. For the aerospace industry, LS-DYNA has ability to simulate bird impact on airframes and engines and turbine rotor burst containment, among others. Additional complexities arise from simulations of these classes since they often require predictions of surface contact and penetration, modeling of loading and material behavior, and accurate failure assessment.

From a hardware and software algorithm perspective, there are roughly three types of LS-DYNA simulation "behavior" to consider: implicit and explicit FEA for structural mechanics, and computational fluid dynamics (CFD) for fluid mechanics. Each discipline and associated algorithms have their inherent complexities with regards to efficiency and parallel performance, and also depending upon modeling parameters.

The range of behavior for the three disciplines that are addressed with LS-DYNA simulations, highlights the importance of a balanced HPC system architecture. For example, implicit FEA for static load conditions requires a fast processor for effective turnaround, in contrast to dynamic response, which requires high rates of memory and I/O bandwidth with processor speed as a secondary concern. In addition, FEA modeling parameters such as the size, the type of elements, and the load condition of interest all affect the execution behavior of implicit and explicit FEA applications.

Explicit FEA benefits from a combination of fast processors for the required element force calculations and a high rate of memory bandwidth necessary for efficient contact resolution that is required for nearly every structural impact simulation. CFD also requires a balance of memory bandwidth and fast processors, but benefits most from parallel scalability. Each discipline has inherent complexities with regard to efficient parallel scaling, depending upon the particular parallel scheme of choice.

Implementations of both shared memory parallel (SMP) and distributed memory parallel (DMP) have been developed for LS-DYNA. The SMP version exhibits moderate parallel efficiency and can be used with SMP computer systems only, while the DMP version, exhibits very good parallel efficiency. This DMP approach is based on domain decomposition with MPI for message passing, and is available for homogenous compute environments such as shared memory parallel systems or clusters.

Most parallel MCAE software employ a similar DMP implementation based on domain decomposition with MPI. This method divides the solution domain into multiple partitions of roughly equal size in terms of required computational work. Each partition is solved on an independent processor, with information transferred between partitions through explicit message passing software (usually MPI) in order to maintain the coherency of the global solution.

LS-DYNA is carefully designed to avoid major sources of parallel inefficiencies, whereby communication overhead is minimized and proper load balance is achieved. Another advantage for LS-DYNA is the use of an SGI-developed MPI that is NUMA-aware and transparent to the user. This MPI further reduces communication overhead when scaling to a large number of processors, which is achieved by a reduction in latency that is improved over public-domain MPI such as MPICH and LAM.

### **SGI IRIX and Linux HPC Server Architectures**

SGI servers are based on a cache-coherent non-uniform memory access (NUMA) multiprocessor architecture that is a breakthrough implementation of conventional shared memory architectures. The SGI NUMA architecture distributes memory to individual processors through a non-blocking interconnect design, in order to reduce latencies that inhibit high bandwidth and scalability. At the same time, all memory is globally addressable, meaning memory is physically distributed but appears logically as a shared resource to the user, to accommodate high-resolution MCAE modeling and simplify MCAE algorithm development.

The motivation for this distributed shared memory NUMA approach evolved at SGI when conventional shared-bus architectures began to exhibit high-latency bottlenecks as the number of processors increased within a single system shared-memory image. During this same time, non-coherent distributed memory architectures started to emerge, but the programming of applications for message passing in such an environment was considered too complex for commercial success.

This SGI NUMA architecture was introduced in the SGI Origin 2000 server in 1995 and later advanced with the SGI® NUMAflex™ modular design concept of the SGI Origin 3000 series. A single image SGI Origin 3900 system offers up to 512 processors and can expand to 1 TB of memory, which is the largest shared memory scalable system available in industry. It is based on an SGI proprietary platform of the IRIX® operating system, the MIPS® microprocessor, and NUMAflex.

Since January 2003, this same NUMAflex architecture of the SGI Origin series has been available in the SGI Altix 3000, an HPC platform based on 64-bit Linux and the Itanium 2 microprocessor from Intel. The SGI Altix 3000 is the first Linux-based HPC system that retains supercomputing features such as global shared-memory for a single image, yet also offers conventional cluster capability. It is an open standards-based platform that combines industry leading HPC system technologies.

The SGI Altix 3000 is built-up from a number of component modules, or bricks, most of which are shared with the Origin 3000 family. The C-brick (compute brick) is the module that customizes the system to a given processor architecture. The SGI Altix 3000 C-brick consists of four Intel Itanium 2 and corresponding memory up to 8GB, and a connection to a portion of an I/O subsystem. The hub interface to the C-brick is the distributed memory controller and C-bricks are connected together via the R-brick, which is the NUMAflex router module.

The maximum configuration of a single node SGI Altix 3000 is 64 processors and 512 GB of shared memory within each single Linux OS-image node. These individual 64 processor nodes can be clustered with a choice of scalable interconnect networks, including the proprietary SGI® NUMAlink™ interconnect technology, to much larger system configurations of up to 2,048 processors and a total of 16TB of memory – a configuration that makes up an SGI Altix 3000 supercluster.

The SGI Altix 3000 is binary compatible with the industry-standard 64-bit Linux distribution. In addition, SGI offers differentiated middleware and other functionality to enhance technical HPC workloads in a software module called ProPack -- a set of user tools that ride on top of Linux. ProPack is similar to other commercial software packages and is used to boost the performance of Linux and user applications on the SGI Altix 3000, not to alter Linux itself. Such enhancements are commonplace in the industry and are offered by several Linux system providers.

Various other system level tools such as the Intel compilers, the SGI-developed SCSL library, commands runon, dplace, and others are available on the SGI Altix 3000 to help LS-DYNA achieve its performance goals. The general principle of HPC performance tuning for the SGI Altix 3000 is to ensure compiler software pipelining is invoked so that efficient instruction scheduling occurs, which provides good processor, cache, and memory locality of data during execution.

The typical practice of an LS-DYNA user environment is to combine moderate scalability of up to 24 processors for a single job with multi-job throughput, in a mix that provides the greatest overall cost benefits of a particular application. The SGI Altix 3000 combines the HPC technologies that offer industry and manufacturing research organizations a high-availability, non-degrading, and efficient LS-DYNA application environment that ensures turnaround and throughput are delivered in support of hundreds of simultaneous users with a mix of MCAE disciplines.

### **Performance Considerations of LS-DYNA**

Performance and parallel efficiency of any MCAE software has certain algorithm considerations that must be addressed. The fundamental issues behind parallel algorithm design are well understood and described in various research publications. For grid-based problems such as the numerical solution of partial differential equations, there are four main sources of overhead that can degrade ideal parallel performance: 1) non-optimal algorithm overhead, 2) system software overhead, 3) computational load imbalance, and 4) communication overhead.

Parallel efficiency for LS-DYNA is dependent upon among others, MPI latency, which is determined by both the specifics of system architecture and the implementation of MPI for that system. Since system architecture latency is determined by design of a particular interconnect, overall latency improvements can only be made to the MPI implementation. Modifications to the MPI software to ensure "awareness" of a spe-

cific architecture are the only way to reduce the total latency and subsequently the communication overhead. Improvements to architecture-awareness of MPI is a common development for several system vendors including SGI.

Parallel efficiency bottlenecks for LS-DYNA on SGI NUMA systems are identified as MPI latency and non-enforcement of processor-memory affinity (data placement) as the most dominant. The latency problem was solved by invoking SGI NUMA-aware MPT library, which appears in LS-DYNA for both the SGI Origin 3000 and SGI Altix 3000. The data placement concern is related to a feature of the NUMA architecture and is addressed through implementation of the SGI dplace set of tools that appear with both IRIX and Linux on SGI systems.

Specifically for structural FEA simulations, they often contain a mix of materials and finite elements that can exhibit substantial variations in computational expense, which may create load-balance complexities. The ability to efficiently scale to a large number of processors is highly sensitive to load balance quality. For example, the crash worthiness of automotive vehicles exhibit these characteristics.

An additional consideration for improved load balance is efficient use of the domain decomposer. The automotive model in Figure 1. illustrates the results of a geometric-based RCB decomposition with use of (left) default parameters and (right) use of an improved scheme that exhibits better overall performance.



**Figure 1. LS-DYNA Decomposition Results for Default (left) and Modified (right) RCB**

For this particular model and partitioning strategy, the modified RCB achieved an improvement of 37% over the default RCB partitioning for 8 processors, and this advantage grew to 55% for 16 processors. This is due to the fact that the modified RCB provides better data partitioning with a more even distribution of processor workloads and less communication between processors.

As a result of collaboration with joint technologies of HPC systems from SGI and FEA software from LSTC, LS-DYNA when executed on an SGI Origin 3000 series system, often exhibits linear scalability as high as 96 processors for various industrial-sized automotive vehicle models. This substantial savings in simulation turnaround time allows additional studies to be performed towards optimization of a vehicle's crash-worthiness and safety.

Next parallel scalability performance is presented in Figure 2. to demonstrate the abilities of a moderately configured SGI Altix 3000 system. The system contained 32 Itanium 2 processors from Intel, in globally shared memory arrangement with 32 GB

of memory. The model for this test was an automotive vehicle, a Dodge Neon configuration, with roughly 535,000 elements, with a simulation time of 30 milliseconds.

The use of 30-milliseconds ensures that the model considers the effect of contact in its ability to scale. During the first 5 milliseconds of simulation there are essentially no contacts, which would provide a true test of parallel scalability for a fully converged simulation.

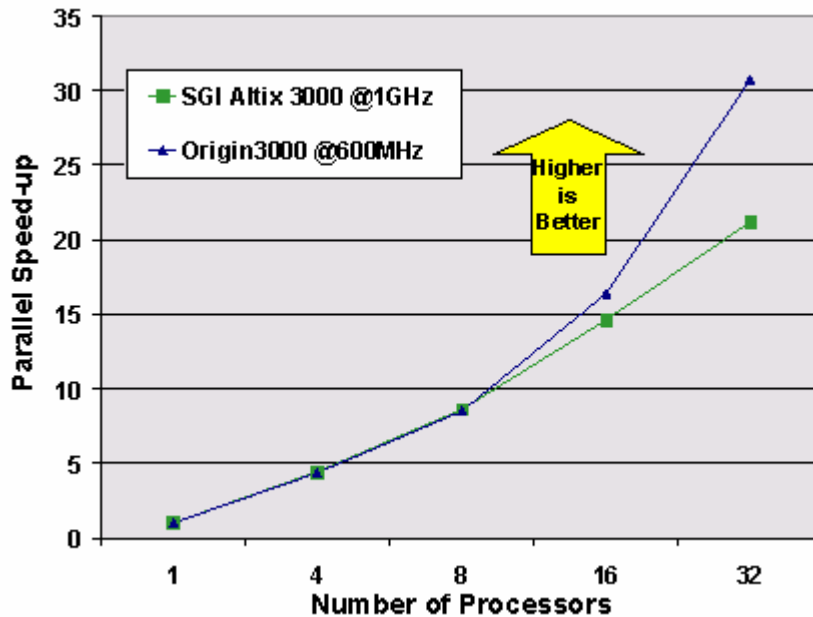


Figure 2. LS-DYNA Scalability of the Dodge Neon Case of 534K Elements

Evaluation of an vehicle's crashworthiness is currently the fastest growing application for MCAE simulation in the automotive industry. Proper crash management and occupant safety is a mandate by local governments, but its also viewed as a competitive advantage by many automotive developers. Performance improvements continue with LS-DYNA such that parallel scalability keeps pace with growing demand.

Similarly, an aerospace application for design of gas turbine engines for aircraft, has utilized the parallel scalability of LS-DYNA on a 64-processor Origin 3000 series system to reduce from 156 hours to 2.5 hours the time it takes to develop a 500,000-element model for blade-out simulation. Parallel efficiency for SGI Altix 3000 for this same model is only slightly lower than Origin 3000, demonstrating the benefit of the SGI NUMAflex architecture for Linux and Intel Itanium 2 microprocessors.

Additional developments between LSTC and SGI applications engineering include an enhanced I/O scheme that significantly improves overall model turnaround in a mix of LS-DYNA jobs in an HPC production environment. This development is particularly important in production environments that include other applications and disciplines such as NVH and CFD that might request similar HPC resources as LS-DYNA during a multi-job throughput workload.

### **Summary and Conclusions**

A discussion was provided on the HPC technology requirements of LS-DYNA applications, including characterizations of the performance behavior typical of LS-DYNA simulations on SGI Origin 3000 and SGI Altix 3000 servers. Effective implementation of highly parallel LS-DYNA simulations must consider a number of features such as parallel algorithm design, system software performance issues, and hardware communication architectures.

Development of increased parallel capability will continue on both application software and hardware fronts to enable FEA modeling at increasingly higher resolutions. Examples of LS-DYNA simulations demonstrate the possibilities for highly efficient parallel scaling on the SGI NUMA systems.

LSTC and SGI will continue to develop software and hardware performance improvements, enhanced features and capabilities, and greater parallel scalability to accelerate the overall solution process of LS-DYNA simulations. This alliance has improved FEA modeling practices in research and industry on a global basis and will continue to provide advancements for a complete range of engineering applications.

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